

PMIN - A Recursive Quadratic Programming Variable Metric Algorithm for Minimax Optimization

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PMIN - A Recursive Quadratic Quadratic Quadratic Quadratic Quadratic Quadratic Quadratic Quadratic Quadratic Q Programming Variable Metric Algorithm for Minimax Optimization

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PMIN - A Recursive Quadratic Programming Variable Metric Algorithm for Minimax Optimization

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Abstract. We present FORTRAN subroutines for nonlinear minimax optimization with simple bounds and general linear constraints based on a recursive quadratic pro gramming variable metric algorithm-

Categories and Subject Descriptors:

General Terms: Algorithms

Additional Key Words and Phrases: Minimax optimization, discrete Chebyshev approximation, recursive quadratic programming methods, variable metric methods, general linear constraints

The double-precision FORTRAN 77 basic subroutine PMIN is designed to find a close approximation to a local minimum of a special ob jective function

$$
F(x) = \max_{1 \le i \le n_a} f_i(x)
$$

(minimax) with simple bounds on variables and general linear constraints. Here $x \in R^n$ is a vector of *n* variables and $f_i: R^n \to R, 1 \le i \le m$, are twice continuously differentiable functions. Simple bounds are assumed in the form $(I^-,I^-$ correspond to $\hspace{0.1mm}$ the arays IV-disc in Section , \mathbf{I}

$$
x_i
$$
 – unbounded , $I_i^x = 0$,
\n $x_i^l \le x_i$, , $I_i^x = 1$,
\n $x_i \le x_i^u$, , $I_i^x = 2$,
\n $x_i^l \le x_i \le x_i^u$, , $I_i^x = 3$,
\n $x_i = x_i^l = x_i^u$, , $I_i^x = 5$,

where $1 \leq i \leq n$. General linear constraints are assumed in the form

 $a_i^T x$ – unbounded, $I_i^c = 0$,

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$$
c_i^l \le a_i^T x \quad , \quad I_i^c = 1,
$$

\n
$$
a_i^T x \le c_i^u \quad , \quad I_i^c = 2,
$$

\n
$$
c_i^l \le a_i^T x \le c_i^u \quad , \quad I_i^c = 3,
$$

\n
$$
a_i^T x = c_i^l = c_i^u \quad , \quad I_i^c = 5,
$$

where $1 \leq i \leq n_c$ and n_c is a number of general linear constraints. To simplify user's work three additional easy to use subroutines are added-to substitute and addedsubroutine PMIN

PMINU  unconstrained minimax optimization

PMINS  minimax optimization with simple bounds

PMINL  minimax optimization with simple bounds and general linear constraints

All subroutines contain a description of formal parameters and extensive comments-Furthermore, two test programs TMINU and TMINL are included, which contain 7 and it is the problems see the server is the programs serves for the serves for the server is the server of th subroutines, verify their correctness and demonstrate their efficiency.

To simplify the description of the method, we will consider a simpler problem written in the following compact form

$$
x^* = \arg\min_{x \in L^n} (\max_{i \in M_1} f_i(x)),\tag{1}
$$

where

$$
L^n = \{x \in R^n : a_i^T x \le b_i, i \in M_2\}
$$

with $M_1 \cap M_2 = \emptyset$. It is clear that the application of the method described below to the general problem stated in the previous section is straightforward, but this requires to consider each type of constraint separately as it is realized in the subroutine PMIN-

- Recursive quadratic programming variable metric method for nonlinear minimax optimization

If we introduce a new variable z, then the problem (1) can be reformulated as a nonlinear programming problem

$$
(x^*, z^*) = \arg\min_{(x,z)\in N^{n+1}} z,
$$
\n(2)

where

$$
N^{n+1} = \{(x, z) \in R^{n+1} : f_i(x) \le e_i z, i \in M_1 \cup M_2\}
$$

with $e_i = 1$ for $i \in M_1$ and $e_i = 0$, $f_i(x) = a_i^T x - b_i$ for $i \in M_2$. This nonlinear programming problem can be solved by a recursive quadratic programming method that uses a quadratic approximation of the Lagrangian function and a linear approximation of the constraints in each iteration. Let $x^k \in R^n$ be a current approximation to the minimum $x^{\scriptscriptstyle \top}$. Then the resulting quadratic programming subproblem has the form

$$
(d^k, z^k) = \arg\min_{(d,z)\in L_k^{n+1}} \left(\frac{1}{2}d^T G^k d + z\right),\tag{3}
$$

where G^k is an approximation of the Hessian matrix of the Lagrangian function and

$$
L_k^{n+1} = \{(d, z) \in R^{n+1} : f_i^k + (a_i^k)^T d \le e_i z, i \in M_1 \cup M_2\}
$$

with $e_i = 1, a_i^k = \nabla f_i(x^k)$ for $i \in M_1$ and $e_i = 0, a_i^k = a_i$ for $i \in M_2$. The solution of the quadratic programming subproblem (3) has to satisfy the Karush-Kuhn-Tucker conditions

$$
d^{k} = -H^{k}g^{k}
$$

\n
$$
e^{T}u^{k} = 1,
$$

\n
$$
u^{k} \geq 0,
$$

\n
$$
v^{k} \geq 0,
$$

\n
$$
v^{k})^{T}u^{k} = 0,
$$

where u is the vector of Lagrange multipliers, $H^+ = (G^+)$, $A^+ = [a_1^+, \ldots, a_m^+]$, $e =$ $[e_1,\ldots,e_m]$, $f^*= [f_1,\ldots,f_m]$, $g^*=A^*u^*$ is the gradient of the Lagrangian function and $v^k = z^k e - f^k - (A^k)^T d^k$ is the vector of constraint violations. Note that if $g^k = 0$, then we obtain the Karush-Kuhn-Tucker conditions for the nonlinear programming , problem , we have the minimax conditions the minimax , we have the conditions the conditions of the condition $||g^k||_{\infty} \leq TOLG$ is used in the subroutine PMIN as the basic stopping criterion (when it is fulfilled then $ITERM=4$).

The direction vector $d^k \in \mathbb{R}^n$ obtained as the solution to the quadratic programming subproblem (3) is used for the definition of the new approximation x^{k+1} to the minimum x – by the formula –

$$
x^{k+1} = x^k + \alpha^k d^k,
$$

where $0 < \alpha^{k} < 1$ is a steplength, which is chosen in such a way that

 $\left(\cdot\right)$

$$
F(x^k + \alpha^k d^k) - F(x^k) \leq \underline{\varepsilon} \alpha^k (d^k)^T g^k,
$$

where is a tolerance for function decrease in the line search parameter \mathcal{C} TOLS in the subroutine P MIN). The steplength α^+ is chosen iteratively either by the bisection (MES=1) or by two point quadratic interpolation (MES=2) or by three point quadratic interpolation (MES=3) or by three point cubic interpolation (MES=4) (MES is a parameter of the subroutine PMIN).

Having the new approximation x^{\ldots} to the minimum x^{\ldots} , we we can compute the new computer the new computer the new computer terms of the new computer of the new matrix $A^{k+1} = [a_1^{k+1}, \ldots, a_m^{k+1}]$, where $a_i^{k+1} = \nabla f_i(x^{k+1})$ for $i \in M_1$ and $a_i^{k+1} = a_i$ for $i \in M_2$. If we denote $s^k = x^{k+1} - x^k$ and $y^k = A^{k+1}u^k - A^k u^k = A^{k+1}u^k - g^k$, then the BFGS $([1], [2], [4], [8])$ method consists in the following update

$$
G^{k+1} = \frac{1}{\gamma^k} (G^k + \gamma^k \frac{y^k (y^k)^T}{(s^k)^T y^k} - \frac{G^k s^k (G^k s^k)^T}{(s^k)^T G^k s^k}) = \frac{1}{\gamma^k} (G^k + \gamma^k \frac{y^k (y^k)^T}{(s^k)^T y^k} + \alpha^k \frac{g^k (g^k)^T}{(s^k)^T g^k}),
$$

where γ^+ $>$ 0 is a self scaling parameter. This parameter is usually equal to one $$ with the exception of the first iteration (or iteration after the restart) where either $\gamma^k = 1$ if MET=1 or $\gamma^k = (s^k)^T G^k s^k / (s^k)^T y^k = -\alpha^k (s^k)^T q^k / (s^k)^T y^k$ if MET=2 (MET is a parameter of the subroutine PMIN- The BFGS method requires the condition (s^+) $y^+ > 0$ to be satisfied, which guaralities a positive definiteness of the matrix $\rm G^{++}$. Unfortunately, this condition does not hold in the minimax optimization case automatically. Therefore, we set $G^{k+1} = G^k$ whenever $(s^k)^T y^k \leq 0$.

- Dual range space method for a special quadratic program ming subproblem

Consider a quadratic programming problem in which we seek a pair $(d^*, z^*) \in R^{n+1}$ in such a way that

$$
\phi(d^*, z^*) = \min_{(d,z)\in L^{n+1}} \phi(s, z),
$$
\n(4)

where

$$
\phi(s,z) = \frac{1}{2}d^T G d + z
$$

and

$$
L^{n+1} = \{(d, z) \in R^{n+1} : f_i + a_i^T d \le e_i z, i \in M_1 \cup M_2\}
$$

see is fact that the matrix \sim matrix \sim is positive decreased that the the problems is \sim is convex and we can apply the duality theory to obtain a dual quadratic programming problem which consists in seeking a vector $u^* \in \mathbb{R}^m$ (vector of Lagrange multipliers of (4) so that

$$
\psi(u^*) = \min_{u \in L^m} \psi(u),\tag{5}
$$

where

$$
\psi(u) = \frac{1}{2}u^T A^T H A u - f^T u
$$

 \cdots

and

$$
L^m = \{ u \in R^m : e^T u = 1, u \ge 0 \}.
$$

Here $H = G^{-1}, H = [a_1, \ldots, a_m], f = [f_1, \ldots, f_m]^T, e = [e_1, \ldots, e_m]^T$, where $f_i = f_i(x)$, $e_i = 1$ for $i \in M_1$ and $f_i = a_i^T x - b_i$, $e_i = 0$ for $i \in M_2$. The solution of (4) can be obtained from the solution of (5) by the formulas

$$
d^* = -H A u^* \tag{6}
$$

and

$$
z^* = f^T u^* - (u^*)^T A^T H A u^*.
$$
\n⁽⁷⁾

The solution u° of (5) is the optimum vector of Lagrange multipliers for (4). Since the problem (5) is convex, u^- is its solution if and only if the Karush-Kuhn-Tucker conditions are valid in the valid in the valid in the valid is a set of and only if and only if and only if an

$$
e^T u^\star = 1, \quad u^\star \ge 0 \tag{8}
$$

and a number z^- exists in such a way that \qquad

$$
v^* = A^T H A u^* - f + z^* e \ge 0, \quad (v^*)^T u^* = 0.
$$
 (9)

vector v^- is the vector of Lagrange multipliers of the problem (5). Conditions (0) and (9) imply that z^- in (9) is identical with z^- in (1). This in turn implies that v^- is, at the same time, the vector of constraint values of the problem (4) .

Consider any subset $I \subset M = M_1 \cup M_2$ and denote the vectors of elements u_i, f_i, \ldots $e_i, i \in I$ by u, f, e , respectively. Similarly, let A be the matrix of columns $a_i, i \in I$. To connect two separate cases which can occur in an investigation of a dual range space method together, we introduce and intrinsical parameters, μ we have denote

$$
\tilde{A} = \left[\begin{array}{c} A \\ -e^T \end{array} \right], \quad \tilde{H} = \left[\begin{array}{cc} H & 0 \\ 0 & \eta \end{array} \right].
$$

We will suppose that the subset $I \subset M = M_1 \cup M_2$ was chosen in such a way that the columns of A are linearly independent-

 μ $I = I$ were the set of active constraints of the problem (4) , then we could compute the qual variables z° and u° from (8)-(9). Unfortunately, this set is not known a priory. Therefore, we start with the set $I = \{k\}$, where $k \in M_1$ is arbitrary. Then $z = f_k - a_k^THa_k$ and $u = [1]$. Suppose that $I \subset M = M_1 \cup M_2$ is a current subset and z u are corresponding dual variables-to-respond in the following way-then was-First we compute the direction vector $d = -H Au$ and the value of the most violated primal constraint

$$
v_k = ze_k - f_k - a_k^T d = \min_{i \in M \setminus I} \{ ze_i - f_i - a_i^T d\}.
$$

If $v_k \geq 0$ then the set of active constraints has been detected and the solutions of (4) and been founded been founded between \mathbb{R} and dual and steplengths

$$
\alpha_P = -\frac{v_k}{\beta_k \gamma_k + \delta_k}
$$

\n
$$
\alpha_D = \frac{u_j}{q_{kj} + \gamma_k p_j} = \min_{i \in \overline{I}} \frac{u_i}{q_{ki} + \gamma_k p_i},
$$

where $p = (A^T H A)^{-1} e$, $q_k = (A^T H A)^{-1} A^T H \tilde{a}_k$, $\beta_k = e_k - e^T q_k$, $\gamma_k = \beta_k / p^T e$, $\delta_k =$ $\tilde{a}_k^T(H-HA(A^THA)^{-1}A^TH)\tilde{a}_k \text{ (with } \tilde{a}_k=[a_k,-e_k]^T) \text{ and } \overline{I}=\{i\in I: q_{ki}+\gamma_kp_i>0\} \; .$

If $\beta_k \gamma_k + \delta_k = 0$, then we set $\alpha_P = \infty$. If $I = \emptyset$, then we set $\alpha_D = \infty$. If simultaneously $\alpha_P = \infty$ and $\alpha_D = \infty$, then the problem has no feasible solution. Otherwise we set $\alpha = \min\{\alpha_P, \alpha_D\}$ and compute $z := z + \alpha \gamma_k$, $u := u - \alpha(q_k + \gamma_k p)$, $u_k := u_k + \alpha$, $v_k := (1 - \alpha/\alpha_P)v_k.$ = min{ α_P, α_D } and compute $z := z + \alpha \gamma_k$, $u := u - \alpha (q_k + \gamma_k p)$, $u_k := u_k + \alpha$,

= $(1 - \alpha/\alpha_P)v_k$.

If $\alpha_P \leq \alpha_D$, then the primal step is realized, i.e. we set $I := I \cup \{k\}$, $u := [u^T, u_k]^T$,

 $e := [e^T, e_k]^T$, $A := [A, a_k]$, $A := [A, \tilde{a}_k]$, recompute $d = -HAu$ and determine a new value of the most violated primal constraint and a new index k .

If $\alpha_P > \alpha_D$, then the dual step is realized, i.e. we set $I := I \setminus \{j\}, u := u^{(j)},$ $e := e^{\sqrt{y}}$, $A := A^{\sqrt{y}}$, $A := A^{\sqrt{y}}$, where the upper index in parentheses denotes an element or column which are deleted. Now, two cases can occur. If $I\cap M_1\neq\emptyset,$ then we recompute the primal and dual steplengths and repeat the process with the same element or column which are deleted. Now, two cases can occur. If $I \cap M_1 \neq \emptyset$, then we recompute the primal and dual steplengths and repeat the process with the same index k. If $I \cap M_1 = \emptyset$, then we compute $z := z - v_k$, $e := [e^T, e_k]^T$, $A := [A, a_k]$, $A := [A, \tilde{a}_k]$, recompute $d = -H A u$ and determine the new value of the most violated primal constraint and the new index k .

In $[5]$, it was proved that the above dual range space finds the solutions of quadratic programming problems (4) and (5) after a finite number of steps.

- Description of the subroutines

In this section we describe all subroutines which can be called from the user's program. In the description of formal parameters we introduce a type of the argument that spec ifies whether the argument must have a value defined on entry to the subroutine (I) , whether it is a value which will be returned (O) , or both (U) , or whether it is an auxiliary value (i.e.) is the argument of the arguments of the the theory of the the theory of the changed on α some circumstances especially if improper input values were given- Besides formal parameters, we can use a COMMON $/STAT/$ block containing statistical information. This block, used in each subroutine has the following form:

COMMON STAT NDECF-NRES-NRED-NREM-NADD-NIT-NFV-NFG-NFH

The arguments have the following meaning-

Argument Type Signi
cance

straint additions during the QP solutions-

sub-outines PMINU PMINS PMINS PMINS PMINS PMINS PMIN

The calling sequences are

CALL PMINUNF-NA-X-AF-IA-RA-IPAR-RPAR-F-GMAX-IEXT-ITERM CALL PMINSNF-NA-NB-X-IX-XL-XU-AF-IA-RA-IPAR-RPAR-F-GMAX- IEXT-ITERM CALL PMINLNF-NA-NB-NC-X-IX-XL-XU-CF-IC-CL-CU-CG-AF-IA-RA-

The arguments have the following meaning.

IPAR-RPAR-F-GMAX-IEXT-ITERM

The subroutines PMINU, PMINS, PMINL require the user supplied subroutines FUN and DER that define the values and the gradients of the functions in the minimax criterion and have the form

 $\mathbf{S} = \mathbf{S} \mathbf{S}$

SUBROUTINE DERNF-KA-X-GA

The arguments of user supplied subroutines have the following meaning-

ITERM if the method failed-

Argument Type Signi
cance

3.2 Subroutine PMIN

This general subroutine is called from all the subroutines described in Section - \mathbf{f} calling sequence is

NA-IA-IA-IA-IA-IA-IA-IA-IA-IA-I GA-AG-IAA-AR-AZ-G-H-S-XO-GO-TOLX-TOLF-TOLB-TOLG-TOLD-TOLS-XMAX-GMAX-F-IEXT-MET-MES-MIT-MFV-IPRNT-ITERM

The arguments NF, NA, NB, NC, X, IX, XL, XU, CF, IC, CL, CU, CG, AF, GMAX, F, IEXT, ITERM have the same meaning as in Section -- Other arguments have the following meaning.

Argument Type Signi
cance

the subroutine PMIN has a modular structure- which it is modelling list contains its most contains important subroutines-

- PDDXQ1 Determination of the descent direction using quadratic programming subroutine-
- PLQDF1 Dual range space method for solving the quadratic programming problem with linear constraints (see $[5]$).
- Line search using only function values. PS0LA2
- PUDBG1 The BFGS variable metric update applied to the Choleski decomposition of the approximate Hessian matrix-

The subroutine PMIN requires the user supplied subroutines FUN and DER- User supplied subroutines FUN and DER are described in Section --

3.3 Subroutine PLQDF1

Since the dual range space method for special quadratic programming subproblems arising in nonlinear minimax optimization can be used separately in many applications e-g- in bundletype methods for nonsmooth optimization we describe the subroutine . Play in the calling sequence is the calling sequence in the calling sequence is the calling sequence in the calling sequence

CALL PLQDF NF-NA-NC-X-IX-XL-XU-AF-AFD-IA-IAA-AG-AR-AZ-

- CF-IC-CL-CU-CG-G-H-S-MFP-KBF-KBC-IDECF-ETA-ETA -ETA-
- EPS
-EPS-XNORM-UMAX-GMAX-N-ITERQ

The arguments NF, NA, NC, X, IX, XL, XU, AF, CF, IC, CL, CU, CG have the same meaning as in Section - only with the dierence that the dierence that the arguments X and X and X e-community to the second control to the second control on entry to play and the second changed-complete the s The arguments AFD IA IAA AG AR AZ have the same meaning as in Section - only with the difference that the arguments AFD III, which are of the the type () it are control values can be used subsequently- Other arguments have the following meaning-

ITERQ $= -1$: an arbitrary feasible point does not exist. ITERQ $= -2$: the optimum feasible point does not exist.

3.4 Form of printed results

The form of printed results is specified by the parameter IPRNT as is described above. Here we demonstrate individual forms of printed results by the simple use of the pro \mathcal{L} the printed results will have the form

 $NIT = 16$ $NFV = 18$ $NFG = 17$ $F =$ $.50694800D+00$ $G= .2872D-06$ ITERM= 4

If we set $IPRNT = 1$, then the printed results will have the form

```
EXIT FROM PMIN :
```

```
X= .5000000D+00 .5000000D+00 .5000000D+00 .5000000D+00 .500000D+00
   .5000000D+00 .5000000D+00 .5000000D+00 .5000000D+00 .500000D+00
  -0.4166693D+00 -0.4166693D+00 -0.4166693D+00 -0.4166693D+00
```
If we set IPRNT=2, then the printed results will have the form

```
ENTRY TO PMIN
```


If we set IPRNT $= 2$, then the printed results will have the form

4. Verification of the subroutines

In this section we introduce the main programs TMINU and TMINL which serve as demonstration, verification and testing of the subroutines PMINU and PMINL.

Program TMINU

The following main program demonstrates the usage of the subroutine PMINU-

```
\overline{C}C TEST PROGRAM FOR THE SUBROUTINE PMINU
\overline{C}INTEGER NEGERIA INTEGERIA ILE ENTERET ENTERET ELEMENT ENTERET ELECTRICIAN ELEC
           REAL X-
AF
-
RA-
RPAR
-
F-
GMAX
          REAL*8 FMIN
           INTEGER NAL-
NEXT-
IERR-
I
          COMMON / PROB/ NEXT
           INTEGER NDECF-
NRES-
NRED-
NREM-
NADD-
NIT-
NFV-
NFG-
NFH
```
COMMON STAT NDECF-NRES-NRED-NREM-NADD-NIT-NFV-NFG-NFH

```
\mathcal{C}C LOOP FOR 7 TEST PROBLEMS
       <u>-- - ----- - , .</u>
C
C CHOICE OF INTEGER AND REAL PARAMETERS
       DO 	 I	-

       IPAR(I)=01 CONTINUE
       DO 
 I	-

       RPAR(I)=0.0D 0
       IPAR(5)=1\mathbf CC PROBLEM DIMENSION
\mathsf{C}NF = 20NA = 30C INITIATION OF X AND CHOICE OF RPAR(7)
       CALL TIUDNF-
NA-
NAL-
X-
FMIN-
RPAR
-
NEXT-
IEXT-
IERR
       IF (IERR.NE.O) GO TO 3
\overline{C}\mathsf{C}CALL PMINUNF-
NA-
X-
AF-
IA-
RA-
IPAR-
RPAR-
F-
GMAX-
IEXT-
ITERM
     3 CONTINUE
       STOP
\mathsf{C}C USER SUPPLIED SUBROUTINE CALCULATION OF FA
       SUBROUTINE FUNNF-
KA-
X-
FA
       INTEGER NEGERIA DEL GERMANIA DE
       REAL X-
FA
       INTEGER NEXT
       COMMON / PROB/ NEXT
\mathsf{C}C FUNCTION EVALUATION
       CALL TAFUNF-
KA-
X-
FA-
NEXT
```

```
16
```

```
RETURN
      END
C USER SUPPLIED SUBROUTINE CALCULATION OF GA
\overline{C}SUBROUTING THE RESIDENCE OF LAND CONTINUES.
       INTEGER NF-
KA
       REAL X-
GA
      INTEGER NEXT
       COMMON / PROB/ NEXT
\overline{C}C GRADIENT EVALUATION
\overline{C}CALL TAGUNF-
KA-
X-
GA-
NEXT
      RETURN
      END
```
This matrix matrix \mathbf{I} is matrix to the subset of \mathbf{I} tion and TAGU subgradient evaluation containing standard test problems with at most variables which were taken from the UFO system - The results obtained by this main program have the following form-

The rows corresponding to individual test problems contain the number of iterations NIT, the number of function evaluations NFV, the number of gradient evaluations NFG, the final value of the objective function F, the value of the criterion for the termination G and the cause of termination ITERM-

4.2 Program TMINL

The following main program demonstrates the usage of the subroutine PMINL-

```
\overline{C}\mathsf{C}TEST PROGRAM FOR THE SUBROUTINE PMINL
                 INTEGER NF-
NA-
NB-
NC-
IX-
IC	-
IA
-
IEXT-
IPAR-
ITERM
                 real contract and the contract of the contract
               ray the contract of the contra
```

```
REAL*8 FMIN
       INTEGER NAL-
NCL-
NEXT-
IERR-
I
      COMMON / PROB/ NEXT
       INTEGER NDECF-
NRES-
NRED-
NREM-
NADD-
NIT-
NFV-
NFG-
NFH
       COMMON STAT NDECF-
NRES-
NRED-
NREM-
NADD-
NIT-
NFV-
NFG-
NFH
C
C LOOP FOR 6 TEST PROBLEMS
C
       do next and next and
C
C CHOICE OF INTEGER AND REAL PARAMETERS
\overline{C}DO 	 I	-

      IPAR(I)=01 CONTINUE
       \overline{\phantom{a}}RPAR(I)=0.0D 0
      IPAR(5)=1C PROBLEM DIMENSION
      NA = 165NC=10C
C INITIATION OF X AND CHOICE OF RPAR(7)
\overline{C}NAL-ILI ILI ILI ILI ILI ILI ILI ILI
      NEXT-
IEXT-
IERR
      IF (IERR.NE.O) GO TO 3
C
C SOLUTION
       CALL PMINLNF-
NA-
NB-
NC-
X-
IX-
XL-
XU-
CF-
IC-
CL-
CU-
CG-
AF-
IA-
RA-
IPAR-
      RPAR-
F-
GMAX-
IEXT-
ITERM
    3 CONTINUE
      STOP
      END
C
C USER SUPPLIED SUBROUTINE CALCULATION OF FA
C
```

```
\mathcal{S} = \{x_1, x_2, \ldots, x_n\}INTEGER NF-
KA
       REAL X-
FA
       INTEGER NEXT
       COMMON / PROB/ NEXT
C FUNCTION EVALUATION
       NF-RESERVED IN THE RESERVED OF THE RESERVED OF
       RETURN
       END
\mathsf{C}C USER SUPPLIED SUBROUTINE CALCULATION OF GA
Ċ.
       SUBROUTING THE RESIDENCE OF REAL PROPERTY.
       \overline{I}REAL X-
GA
       INTEGER NEXT
       COMMON / PROB/ NEXT
C GRADIENT EVALUATION
       CALL TAGU

NF-
KA-
X-
GA-
NEXT
       RETURN
       END
```
This main program uses subroutines TIUD22 (initiation), TAFU22 (function evaluation), TAGU22 (subgradient evaluation) containing 6 standard test problems with at most variables which were taken from the UFO system - The results obtained by this main program have the following form-

The rows corresponding to individual test problems contain the number of iterations NIT, the number of function evaluations NFV, the number of gradient evaluations NFG, the final value of the objective function F, the value of the criterion for the termination G and the cause of termination ITERM-

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